

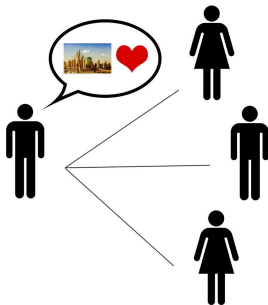
Strategic Disclosure of Opinions on a Social Network

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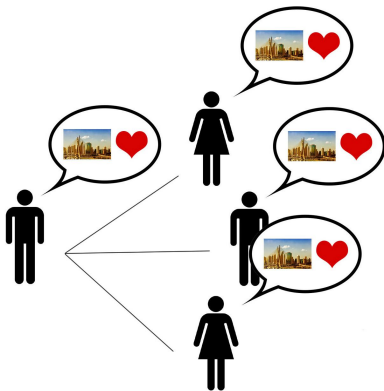


Introductory Example



"São Paulo is the best city in the world!"

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Talk Overview

1. Opinion Diffusion Process
2. Games of Influence
3. Computational Complexity
4. Conclusion and Future Work

Formal Framework

Opinions, influence powers, states

- N is a set of n **agents**
- I is a set of m **issues**

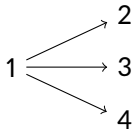
- $B_i : I \rightarrow \{0, 1\}$ is the **opinion** of agent i
 $B_i(p) = 1$ iff i believes that p
- $V_i : I \rightarrow \{0, 1\}$ is the **influence power** (visibility) of agent i
 $V_i(p) = 1$ iff i influences others about her opinion on p

- A **state** consists of all opinions and visibilities of all the agents

Formal Framework

Influence network, update via aggregation

- An **influence network** is a directed irreflexive graph $E \subseteq N \times N$
 $(i, j) \in E$ iff agent i influences agent j



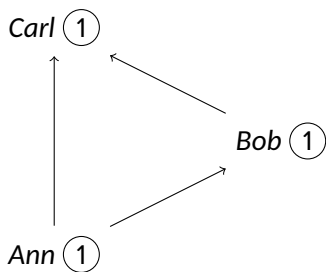
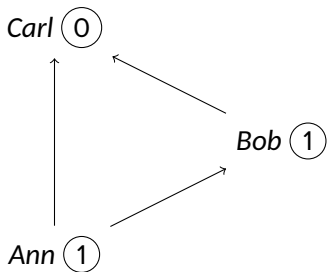
The **opinion update** is a two-step process:

1. Agents decide how to use their influence power on issues
2. Agents update opinions via an aggregation procedure

Unanimous Aggregation

We focus on the **unanimous aggregation procedure**

Agent i updates her opinion on p iff all i 's influencers using their influence are **unanimous** (else, she keeps her current opinion)



Games of Influence

An **influence game** is a tuple $IG = \langle N, I, E, \{F_i\}_{i \in N}, S_0, \{\gamma_i\}_{i \in N} \rangle$

- S_0 is the **initial state**
- γ_i is the **individual goal of agent i**

Goals are expressed in Linear Temporal Logic, with atoms:

- $op(i,p)$ " i believes that p "
- $vis(i,p)$ " i uses her influence power on p "

$$\text{Influence}(i, C, J) := \diamond \square \bigwedge_{p \in J} ((op(i, p) \rightarrow \bigcirc pcon(C, p)) \wedge (\neg op(i, p) \rightarrow \bigcirc ncon(C, p))).$$

Strategies and Solution Concepts

Agents have **actions** $\text{reveal}(J)$ and $\text{hide}(J')$ — use your influence power on the issues in J and not on the issues in J'

We consider two types of strategies

- **Memory-less**: associate action to state
- **Perfect-recall**: associate action to finite sequence of states

Solution concepts (for this talk)

Weakly dominant strategy: agent doesn't gain with different strategy

Nash equilibrium: no agent gains by changing (alone) her strategy

Game Theoretic Results

Memory-less strategies

Always using your influence power is not necessarily a dominant strategy for the Influence goal.

	A	→ B	→ C	→ D
S_0 :	0	1	0	1
S_1 :	0	0	1	0
S_2 :	0	0	0	1

- $\gamma_B = \text{Influence}(B, \{D\}, \{p\})$
 - **B**: always use influence power over p
 - C: use influence power over p *unless* A, B and C agree on p
- ⇒ What if **B** does *not* use her influence power over p in S_0 ?

Computational Complexity Results

Memory-less strategies

M-Nash: Given IG and strategy profile, is it a NE of IG?

Theorem

M-Nash is in P-SPACE for memory-less strategies.

- Encoding of unanimity rule and strategies as LTL formulas
- Validity checking for LTL is in P-SPACE

Computational Complexity Results

Perfect-recall strategies

E-Nash: Given IG, is there a NE of IG?

U-Nash: Given IG, is there a unique NE of IG?

Theorem

Both problems are in 3-EXPTIME for perfect-recall strategies.

- Translation into Graded Strategy Logic formulas
- Model checking of these formulas over a corresponding CGS

B. Aminof, V. Malvone, A. Murano, and S. Rubin. Graded Strategy Logic: Reasoning about Uniqueness of Nash Equilibria (AAMAS-2016).

Conclusion and Future Work

- Extended POD model with minimal strategic element
 - Complex setting (both game-theoretically & computationally)
 - Interesting connection with iterated boolean games
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- Allow agents more actions (e.g., lying)
 - Study different aggregation procedures